Memory Safety

Question 1  Software Vulnerabilities

For the following code, assume an attacker can control the value of basket passed into search_basket. The value of n is constrained to correctly reflect the number of cats in basket.

The code includes several security vulnerabilities. Circle three such vulnerabilities in the code and briefly explain each of the three on the next page.

```c
struct cat {
    char name[1024];
    int age;
};

/* Searches through a basket of cats of size at most 32.
   Returns the number of kittens or -1 if there are no kittens */
int search_basket(struct cat basket[], size_t n) {
    struct cat kittens[32];
    char kitten_names[1024], cmd[1024];
    int i, total_kittens = 0, names_size = 0;

    if (n > 32) return -1;

    for (i = 0; i <= n; i++) {
        if (basket[i].age < 12) {
            size_t len = strlen(basket[i].name);
            snprintf(kitten_names + names_size, len, "%s ", basket[i].name);
            kittens[total_kittens] = basket[i];
            names_size += len;
            total_kittens += 1;
        }
    }

    if (total_kittens > 5) {
        const char *fmt = "adopt-kittens --num_kittens %d --names %s";
        snprintf(cmd, sizeof(cmd), fmt, total_kittens, kitten_names);
        system(cmd);
    }

    return total_kittens;
}
```

Reminders:

- `snprintf(buf, len, fmt, ...)` works like `printf`, but instead writes to `buf`, and won’t write more than `len - 1` characters. It terminates the characters written with a ‘\0’.
- `system` runs the shell command given by its first argument.
1. Explanation:

________________________________________________________________________
________________________________________________________________________

2. Explanation:

________________________________________________________________________
________________________________________________________________________

3. Explanation:

________________________________________________________________________
________________________________________________________________________

Describe how an attacker could exploit these vulnerabilities to obtain a shell:

________________________________________________________________________
________________________________________________________________________

________________________________________________________________________
Question 2  

C Memory Defenses

Mark the following statements as True or False and justify your solution. Please feel free to discuss with students around you.

1. Stack canaries completely prevent a buffer overflow from overwriting the return instruction pointer.

2. A format-string vulnerability can allow an attacker to overwrite values below the stack pointer.

3. An attacker exploits a buffer overflow to redirect program execution to their input. This attack no longer works if the data execution prevention/executable space protection/NX bit is set.

4. If you have a non-executable stack and heap, buffer overflows are no longer exploitable.

5. If you use a memory-safe language, some buffer overflow attacks are still possible.

6. ASLR, stack canaries, and NX bits all combined are insufficient to prevent exploitation of all buffer overflow attacks.

Short answer!

1. What vulnerability would arise if the canary was above the return address?

2. What vulnerability would arise if the stack canary was between the return address and the saved frame pointer?

3. Assume ASLR is enabled. What vulnerability would arise if the instruction \texttt{jmp *esp} exists in memory?
Question 3  **TCB (Trusted Computing Base)**

In lecture, we discussed the importance of a TCB and the thought that goes into designing it. Answer these following questions about the TCB:

1. What is a TCB?

2. What can we do to reduce the size of the TCB?

3. What components are included in the (physical analog of) TCB for the following security goals:
   
   (a) Preventing break-ins to your apartment

   (b) Locking up your bike

   (c) Preventing people from riding BART for free

   (d) Making sure no explosives are present on an airplane
The following code runs on a 32-bit x86 system. All memory safety defenses are disabled. The compiler does not rearrange stack variables.

```c
void doit(char *s1) {
    char buffer[16];
    strcpy(buffer, s1);
    printf("%s\n%s\n", buffer, s1);
}

int main() {
    char s1[64];
    fgets(s1, sizeof s1, stdin);
    doit(s1);
}
```

(a) Which line contains a memory safety vulnerability? What is the name of the vulnerability present on that line?

(b) Complete the diagram of the stack, right before line 3. Assume normal (non-malicious) program execution. You do not need to write the values on the stack, only the names. There are no extraneous boxes. As in discussion, the bottom of the page represents the lower addresses.

```
compiler padding = 0x00000000
char ________ [64]
--------------------
--------------------
--------------------
char buffer[16]
```
(c) Now we will exploit the program. There is already shellcode at the address 0xbfffdead. Complete the Python script below in order to successfully exploit the program.

NOTE: The Python syntax ’A’ * n indicates that the character ’A’ will be repeated n times. The syntax \xRs indicates a byte with hex value 0xRs.

```python
s1 = 'A' * ___ + '_____________________________________________________' + \
     '_____________________________________________________'

print s1
```